

TDB-ACC-NO: NN77122921

DISCLOSURE TITLE: Technique to Fill a Complex Polygon With Rectangles.
December 1977.

PUBLICATION-DATA: IBM Technical Disclosure Bulletin,
December 1977, US

VOLUME NUMBER: 20

ISSUE NUMBER: 7

PAGE NUMBER: 2921 - 2924

PUBLICATION-DATE: December 1, 1977 (19771201)

CROSS REFERENCE: 0018-8689-20-7-2921

DISCLOSURE TEXT:

4p. A common problem in computer graphics and mask design is

to fill a complex polygon, which may contain numerous concave corners

or holes, with simpler figures, usually rectangles.
The technique

herein described is designed to fill a complex polygon consisting of horizontal and vertical lines with the minimum number of rectangles.

- An economical way of describing a complex polygon is by

subtracting small rectangles from a large rectangle.
For example, as

shown in Fig. 1, the desired resultant polygon is rectangle A minus

rectangles B, C and D.

- All the rectangles can be described by the coordinates of their

lower left corner and upper right corner, respectively.

- A technique for filling a complex polygon with rectangles is as

follows:

1. Read in a list of rectangle descriptors, consisting of the

lower left and upper right corner coordinates. A typical set, where

the large background rectangle is given first, is:

TABLE

RECT.	#	XLL	YLL	XUR	YUR
1		1.	1.	8.	8.
2		2.	2.	4.	5.
3		4.	3.	6.	6.
4		1.	6.	6.	8.

- 2. Form tables of unique values of X-coordinates and unique values of Y-coordinates from the input data. Sort this table in ascending values. Typically for the input data of the above table:

Sorted X-coordinates: 1., 2., 4., 6., 8.

Sorted Y-coordinates: 1., 2., 3., 5., 6., 8.

- 3. Construct a grid wherein the grid coordinates have the

values of the sorted X and Y values, as determined in step 2. A

typical grid is shown in Fig. 2 for the data in the table.

- Henceforth, each grid box can be addressed by the horizontal

and vertical count of box positions from the origin. A typical grid

box is shown in Fig. 3, where I and J are the box counts for

the horizontal and vertical directions, respectively. Assign to each

grid box a variable $G(I,J)$, which is used to indicate whether it is

filled or empty. When $G(I,J)=1$, the box is filled; when $G(I,J)=0$,

the box is empty. Initialize all elements of the array G to 1.

- 4. Skipping the first input rectangle, test the midpoint of

each grid box to determine if it falls inside any of the input

rectangles. Set the corresponding $G(I,J)$ to 0 if the midpoint falls

inside any rectangle except the first. A typical result

corresponding to the data in the table is shown in Fig. 4.

- 5. At this point, the complex polygon has been

filled with

rectangles which are those grid boxes with a $G(I,J)$ of the value 1.

The number of rectangles is not yet minimized.

- 6. The maximum value of I is NI , and the maximum value of J is

NJ . For the situation in Fig. 4, $NI=4$ and $NJ=5$. Take each grid box

one at a time, indexing over I , from 1 to NI , and indexing over J ,

from 1 to NJ . Scan the grid in the positive I direction and the

positive J direction from the current (I,J) position, and determine

how many continuous filled grid boxes extend in each direction. IRX

is the number of continuous grid boxes in the I , or X , direction, and

JRY is the number in the J , or Y , direction.

- In Fig. 4, for grid box $(1,1)$, $IRX=4$ and $JRY=4$.

Take as a

"fill rectangle" the continuous string of grid boxes which have the

greatest length originating at the current grid box.

If $IRX > JRY$,

the string is horizontal. If $JRY > IRX$, the string is vertical.

Store the lower left and upper right coordinates of this rectangle.

Set the value $G(I,J)$ of any grid box included in this rectangle to

zero. Go on to the next grid box with $G(I,J)=1$, and repeat this step

until all grid boxes have been processed or the entire array G is set

to zero. Fig. 5 shows the fill rectangles derived from Fig. 4. Fig.

6 is the dimensionally correct representation of the fill rectangles.

- 7. Test all fill rectangles to determine if any fill rectangle

abuts another which has the same dimension along the abutment. If

this happens, merge them into a single fill rectangle and test this

new rectangle for abutment with any other fill rectangle. Typical

abutments are shown in Fig. 7.

The procedure described above will find the minimum number of fill rectangles for all configurations of complex polygons anticipated. The minimum number will not be found when multiple singularities occur, such as in Fig. 8. However, to improve upon this would require a higher level of testing which may not be worth the expected improvement.

SECURITY: Use, copying and distribution of this data is subject to the restrictions in the Agreement For IBM TDB Database and Related Computer Databases. Unpublished - all rights reserved under the Copyright Laws of the United States. Contains confidential commercial information of IBM exempt from FOIA disclosure per 5 U.S.C. 552(b) (4) and protected under the Trade Secrets Act, 18 U.S.C. 1905.

COPYRIGHT STATEMENT: The text of this article is Copyrighted (c) IBM Corporation 1977. All rights reserved.

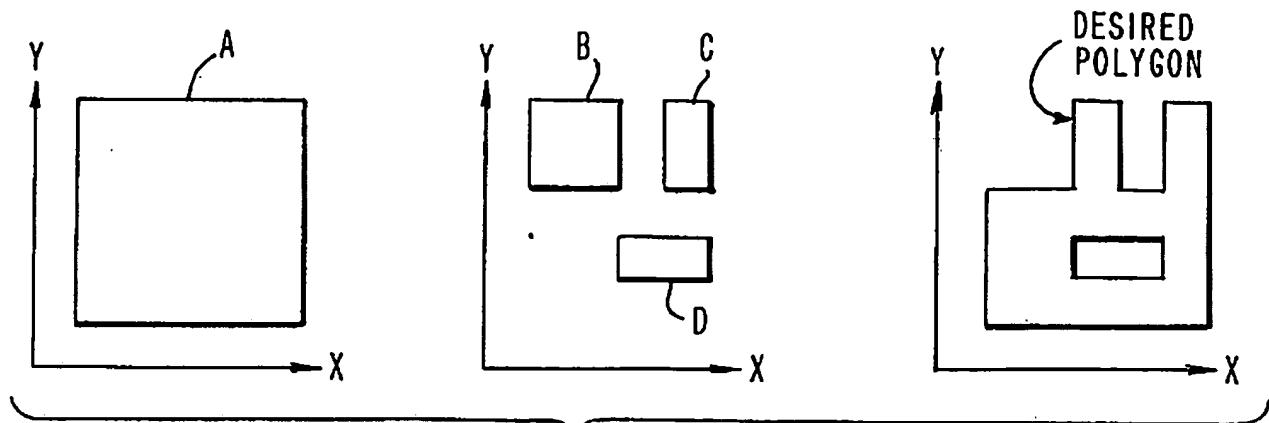
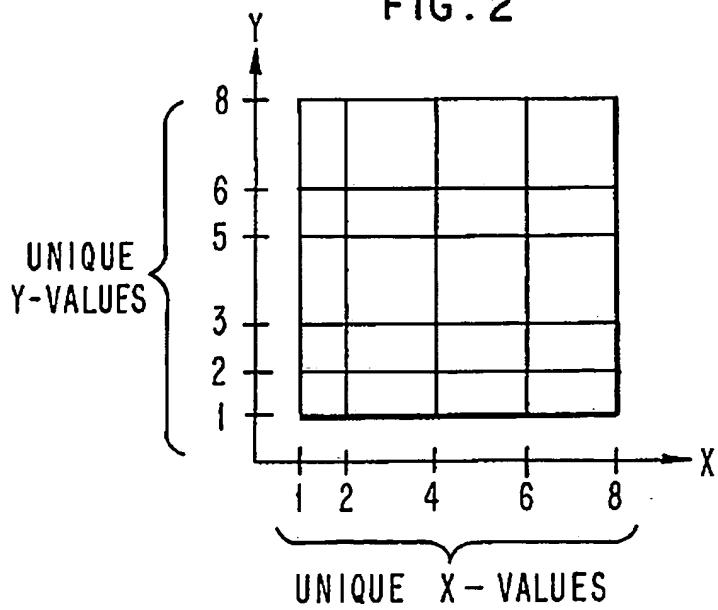


FIG. 1

FIG. 2



GRID BOX (3,4)

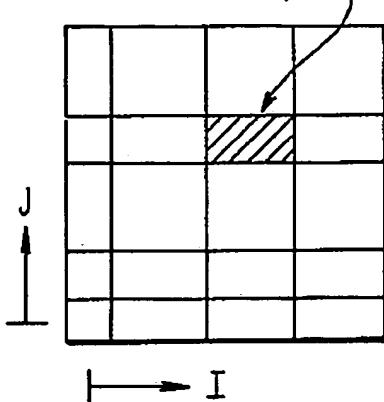


FIG. 3

5	0	0	0	1
4	1	1	0	1
3	1	0	0	1
2	1	0	1	1
1	1	1	1	1
	1	2	3	4

FIG. 4

FIG. 5

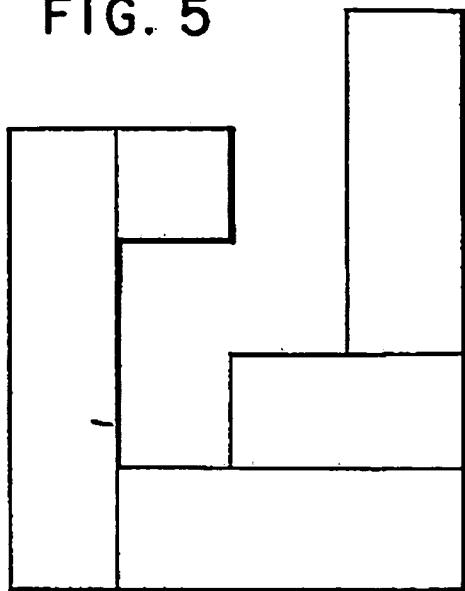


FIG. 6

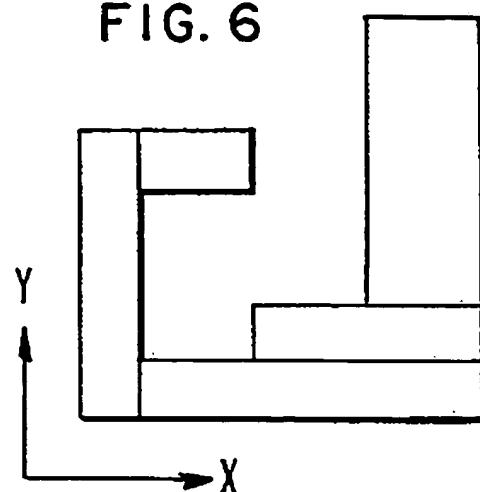


FIG. 7

**MERGER
EXAMPLES**

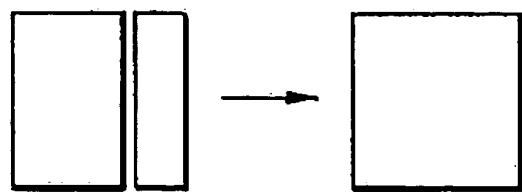


FIG. 8

